Equivalent Stream Reasoning Programs

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- Window function w yields window $w(S, t) \subseteq S$

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- $\Rightarrow$ Do techniques from ASP carry over to LARS?

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- Answer sets characterized by equilibrium models (X, X), where no smaller (X', X) is an SE-model

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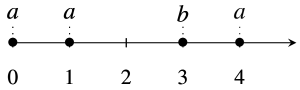

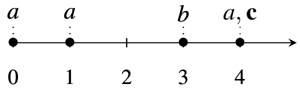

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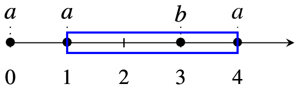

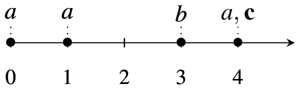

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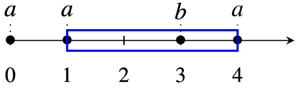

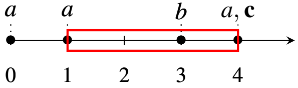

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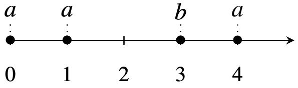

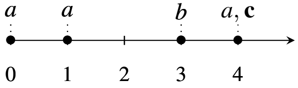

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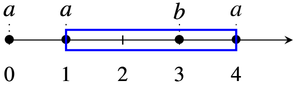

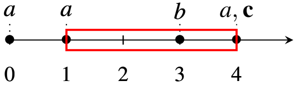

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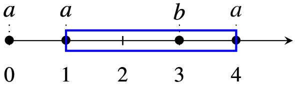

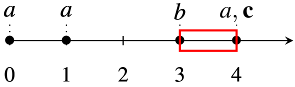

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  - Strong / Uniform / Data Equivalence
- Complexity of deciding eq.: similar to ASP (mostly coNP-c.)