## Decomposition of Distributed Nonmonotonic Multi-Context Systems

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#### **Overview**

Heterogeneous & Nonmonotonic Multi-Context Systems

Decomposition of Nonmonotonic Multi-Context Systems

Experiments

Conclusions

## Multi-Context Systems (MCS)

MCSen introduced by [Giunchiglia and Serafini, 1994]:

- represent inter-contextual information flow
- express reasoning w.r.t. contextual information
- allow decentralized, pointwise information exchange
- monotonic, homogeneous logic
- Framework extended for integrating heterogeneous and nonmonotonic logics [Brewka and Eiter, 2007]

## Syntax of Multi-Context Systems

- multi-context system
  - a collection  $M = (C_1, \ldots, C_n)$  of contexts
- context  $C_i = (L_i, kb_i, br_i)$ 
  - L<sub>i</sub>: a logic
  - *kb<sub>i</sub>*: a knowledge base of logic *L<sub>i</sub>*
  - *br<sub>i</sub>*: a set of bridge rules

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  - *br<sub>i</sub>*: a set of bridge rules
- logic  $L = (\mathbf{KB}_L, \mathbf{BS}_L, \mathbf{ACC}_L)$ 
  - ► **KB**<sub>L</sub>: set of well-formed knowledge bases
  - ► **BS**<sub>L</sub>: is the set of possible belief sets
  - ► ACC<sub>L</sub>: acceptability function KB<sub>L</sub> → 2<sup>BS<sub>L</sub></sup> Which belief sets are accepted by a knowledge base?

## Syntax of Multi-Context Systems (bridge rules)

multi-context system

$$M=(C_1,\ldots,C_n)$$

context

$$C_i = (L_i, kb_i, br_i)$$

Iogic

$$L_i = (\mathbf{KB}_i, \mathbf{BS}_i, \mathbf{ACC}_i)$$

• Bridge rule  $r \in br_i$  of a context  $C_i$ 

$$s \leftarrow (c_1:p_1), \dots, (c_j:p_j),$$
  
 $not(c_{j+1}:p_{j+1}), \dots, not(c_m:p_m)$ 

- $(c_k : p_k)$  looks at belief  $p_k$  in context  $C_{c_k}$
- r is applicable :⇔ positive/negative beliefs are present/absent
- we add the head s to kb<sub>i</sub> if r is applicable

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knowledge base of a context C<sub>i</sub>

 $kb_i \in \mathbf{KB}_i$ 

set of bridge rules br<sub>i</sub> of a context C<sub>i</sub> of form

$$s \leftarrow (c_1:p_1), \ldots, (c_j:p_j), not (c_{j+1}:p_{j+1}), \ldots, not (c_m:p_m)$$

- Contexts C<sub>1</sub>,..., C<sub>n</sub> are knowledge bases with semantics in terms of accepted belief sets
- ►  $S = (S_1, ..., S_n)$  is a belief state of M with each  $S_i \in \mathbf{BS}_i$

## **Semantics of Multi-Context Systems**

multi-context system

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Iogic

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- Equilibrium semantics
  - A belief state  $S = (S_1, \ldots, S_n)$  with  $S_i \in \mathbf{BS}_i$ 
    - ... makes certain bridge rules applicable,
    - $\ldots$  add applicable bridge heads to  $kb_i$
  - $\Rightarrow$  S is an equilibrium : $\Leftrightarrow$

each  $kb_i$  plus acceptable bridge heads from  $br_i$  accepts  $S_i$ 

$$S_i \in \mathbf{ACC}_i(kb_i \cup \{head(r) \mid r \in app(br_i, S)\})$$

#### **The Diamond Example**

 $M = (C_1, C_2, C_3, C_4)$ , were each  $L_i$  of  $C_i$  is an ASP logic



#### Equilibria:

- $\blacktriangleright (\emptyset, \emptyset, \emptyset, \{f\})$
- ▶  $(\emptyset, \{b\}, \{e\}, \{g\})$ 
  - $\blacktriangleright \ (\{a\}, \{b\}, \{c, d\}, \{g\})$

## **Towards Distributed Equilibria Building for MCS**

#### Obstacles:

- abstraction of contexts
- information hiding and security aspects
- lack of system topology
- cycles between contexts

#### We need to capture:

- dependencies between contexts
- representation of partial knowledge
- combination/join of local results

## Import Neighborhood & Closure

Import neighborhood of  $C_k$ 

$$In(k) = \{c_i \mid (c_i : p_i) \in B(r), r \in br_k\}$$



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Import closure IC(k) of  $C_k$  is the smallest set *S* such that (i)  $k \in S$  and (ii) for all  $i \in S$ ,  $In(i) \subseteq S$ .



#### **Partial Belief States and Equilibria**

Let  $M = (C_1, \ldots, C_n)$  be an MCS, and let  $\epsilon \notin \bigcup_{i=1}^n \mathbf{BS}_i$ 

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 $S = (S_1, \ldots, S_n)$  is a partial equilibrium of M w.r.t. a context  $C_k$  iff for  $1 \le i \le n$ ,

- if  $i \in IC(k)$  then  $S_i \in ACC_i(kb_i \cup \{head(r) \mid r \in app(br_i, S)\})$
- otherwise,  $S_i = \epsilon$

Intuitively, partial equilibria w.r.t. a context  $C_k$  cover the reachable contexts of  $C_k$ 













## **Motivation for MCS Decomposition**

Scalability issues with the basic evaluation algorithm DMCS

- unaware of global context dependencies, only know (local) import neighborhood
- a context C<sub>i</sub> returns a possibly huge set of partial belief states, which are the join of neighbor belief states of C<sub>i</sub> plus local belief sets

We address these issues by

- capturing inter-context dependencies (topology)
- providing a decomposition based on biconnected components
- characterizing minimal interface variables in each component
- develop the DMCSOPT algorithm which operates on query plans

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- Problem: How to go home?
- Possible solutions:
  - Car: slower than train
  - Train: should bring some food
- Mr. 3 and Ms. 4 have additional information from Mr. 5 and Ms. 6

- Ms. 4 just got married to Mr. 5.
- Mr. 5 wants his wife to come back as soon as possible.

 $kb_4 = \{car_4 \lor train_4 \leftarrow \}$  $br_4 = \{train_4 \leftarrow (5 : want\_sooner_5)\}$ 

$$kb_{5} = \{want\_sooner_{5} \leftarrow soon_{5}\}$$
  
$$br_{5} = \{soon_{5} \leftarrow (4 : train_{4})\}$$



- Mr. 3 has a daughter, Ms. 6.
- Mr. 3 is responsible for buying provisions, if they go by train.
- If Ms. 6 is sick, then Mr. 3 must attend to her as fast as possible.

$$kb_{3} = \begin{cases} car_{3} \lor train_{3} \leftarrow \\ train_{3} \leftarrow urgent_{3} \\ sandwiches_{3} \lor chocolate\_peanuts_{3} \leftarrow train_{3} \\ coke_{3} \lor juice_{3} \leftarrow train_{3} \end{cases}$$
$$br_{3} = \begin{cases} urgent_{3} \leftarrow (6:sick_{6}) \\ train_{3} \leftarrow (4:train_{4}) \end{cases};$$
$$kb_{6} = \{sick_{6} \lor fit_{6} \leftarrow \}$$

$$br_6 = \emptyset.$$



- Ms. 1 is leader of group.
- ▶ Ms. 1 is allergic to peanuts.
- Mr. 2 wants to get home somehow and doesn't want coke.

$$kb_{1} = \begin{cases} car_{1} \leftarrow not \ train_{1} \\ \perp \leftarrow peanuts_{1} \end{cases}$$
$$br_{1} = \begin{cases} train_{1} \leftarrow (2 : train_{2}), (3 : train_{3}) \\ peanuts_{1} \leftarrow (3 : chocolate\_peanuts_{3}) \end{cases}$$

$$kb_{2} = \{ \perp \leftarrow not \, car_{2}, not \, train_{2} \} \text{ and} \\ br_{2} = \begin{cases} car_{2} \leftarrow (3 : car_{3}), (4 : car_{4}) \\ train_{2} \leftarrow (3 : train_{3}), (4 : train_{4}), \\ not \ (3 : coke_{3}) \end{cases} \end{cases}$$



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One equilibrium is  $S = ({train_1}, {train_2}, {train_3, urgent_3, juice_3, sandwiches_3}, {train_4}, {soon_5, want_sooner_5}, {sick_6})$ 



▶ Ms. 1 decides after gathering information.



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- Mr. 3 and Ms. 4 do not want to bother the others.



- A graph is weakly connected if replacing every directed edge by an undirected edge yields a connected graph.
- ► A vertex *c* of a weakly connected graph *G* is a *cut vertex*, if *G*\*c* is disconnected (3 and 4 are cut vertices)

 Based on cut vertices, we can decompose the MCS into a block tree: provides a "high-level" view of the dependencies (edge partitioning)

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- The block tree of our example is:



- $B_1$  induced by  $\{1, 2, 3, 4\}$
- ▶ *B*<sup>2</sup> induced by {4,5}
- ▶ *B*<sub>3</sub> induced by {3,6}

cycle breaking by creating a spanning tree of a cyclic MCS

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ear decomposition  $P = \langle P_0, \rangle$ 

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ear decomposition  $P = \langle P_0, P_1, \rangle$ 

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cycle breaker edges cb(G, P): remove last edge from each path  $P_i$  in G

Intuition: for a removed edge  $(\ell, t)$ , guess at leaf  $C_{\ell}$  the variables at  $C_{t}$ 

## **Optimization: Avoiding Unnecessary Calls**

*transitive reduction* of a digraph G is the graph  $G^-$  with the smallest set of edges whose transitive closure  $G^+$  equals the one of G



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*transitive reduction* of a digraph G is the graph  $G^-$  with the smallest set of edges whose transitive closure  $G^+$  equals the one of G



Intuitively, the transitive reduction of an acyclic graph is unique, and one can evaluate the contexts using a topological sort of the contexts.



- $B_1$ : acyclic  $\rightarrow$  apply transitive reduction
- B<sub>2</sub>: cyclic → apply ear decomposition, remove last edge from each ear, then apply transitive reduction (already reduced)
- B<sub>3</sub>: acyclic and already reduced

## **Optimization: Minimal Interface**



In a pruned block B', take all variables from

- the minimal interface in B'
- child cut vertices c
- removed edges E

Outcome: query plan for the MCS to restrict calls and partial belief states



- $S = ( \{ train_1 \}, \{ train_2 \}, \{ train_3, urgent_3, juice_3, sandwiches_3 \}, \\ \{ train_4 \}, \{ soon_5, want\_sooner_5 \}, \{ sick_6 \} )$
- $T = ( \{ train_1 \}, \{ train_2 \}, \{ train_3, juice_3, sandwiches_3 \}, \\ \{ train_4 \}, \{ soon_5, want\_sooner_5 \}, \{ fit_6 \} )$
- $U = ( \{ car_1 \}, \{ car_2 \}, \{ car_3 \}, \{ car_4 \}, \emptyset, \{ fit_6 \})$



 $S|_{B_1} = (\{train_1\}, \{train_2\}, \{train_3, juice_3, sandwiches_3\}, \{train_4\}, \emptyset, \emptyset)$  $U|_{B_1} = (\{car_1\}, \{car_2\}, \{car_3\}, \{car_4\}, \emptyset, \emptyset)$ 

## **Experiments**

	п	$A_{\phi}$	$A_{\bowtie}$	$A_{\leftrightarrow}$	$A_{\Sigma}(\sigma)$	# (σ)	$B_{\phi}$	$B_{\bowtie}$	$B_{\leftrightarrow}$	$B_{\Sigma}(\sigma)$	$\#(\sigma)$
D	13	0.9	0.0	0.0	1.0 (0.2)	28 (17.6)	0.8	8.4	0.0	9.4 (5.5)	3136 (3155.8)
	25	11.2	0.5	0.0	12.8 (1.3)	17 (18.9)	_				
	31	51.1	3.7	0.0	59.5 (8.9)	58 (49.7)	—				
R	10	0.1	0.0	0.0	0.1 (0.0)	3.5 (3.4)	0.1	0.0	0.0	0.2 (0.1)	300 (694.5)
	13	0.1	0.0	0.0	0.2 (0.1)	6 (1.2)	0.1	1.5	1.9	3.9 (5.3)	5064 (21523.8)
	301	4.1	0.1	2.1	10.2 (2.2)	8 (4.9)	—				
Ζ	13	0.6	0.1	0.0	0.7 (0.2)	34 (41.8)	5.5	4.2	0.0	11.5 (4.0)	3024 (1286.8)
	151	8.9	22.3	0.4	32.2 (7.3)	33 (28.5)	—				
	301	21.6	99.5	1.7	124.3 (20.6)	22 (41.4)	—				
Η	9	0.2	0.0	0.0	0.2 (0.0)	28 (44.4)	1.1	0.9	0.0	2.0 (1.3)	684 (1308.0)
	101	1.8	0.3	0.3	3.8 (1.0)	48 (76.6)	—				
	301	7.8	2.0	2.4	25.1 (8.7)	38 (34.2)	_				

Table: Runtime for DMCSOPT  $(A_x)$  and DMCS  $(B_x)$ , timeout 180 secs (—) Random instances with *n* contexts and topologies:

Ring  $\xrightarrow{\leftarrow}$ , Zig-zag Diamono . House Timings:

clasp ( $\phi$ ), Belief state combination ( $\bowtie$ ) and transfer ( $\leftrightarrow$ ); No. of partial equilibria: #











#### Conclusions

- MCS is a general framework for integrating diverse formalisms
- First attempt for distributed MCS evaluation
- Initial experiments with a prototype implementation
- Decomposition technique is encouraging: binary tree with n = 600 evaluated in 176secs (# = 4)

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Future work:

- improve scalability
  - approximation semantics
  - syntactic restrictions
  - specialized algorithms for some types of topologies
- dynamic multi-context systems

#### **References I**



Fausto Giunchiglia and Luciano Serafini. Multilanguage hierarchical logics or: How we can do without modal logics.

*Artificial Intelligence*, 65(1):29–70, 1994.