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Business Informatics Group

WIENER WISSENSCHAFTS-
FORSCHUNGS- UND TEGHNOLOCIEFONDS

## Guided Merging of Sequence Diagrams

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## Motivation and Scope



## Motivation and Scope



# Model-Driven Engineering 

## Motivation and Scope



# Model-Driven Engineering 

Model Evolution



## Motivation and Scope



# Model-Driven Engineering 

Model Evolution

Model Versioning

## Motivation and Scope



# Model-Driven Engineering 

Model Evolution

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Model Merging

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# Model-Driven Engineering 

Model Evolution
Model Versioning

Model Merging

Multi-View Modelling

## Example



## Example



## Example



## Example



## Example



## Example



## Example



## Example


$S_{1}$ ?


## Example


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## $S_{1}$ ?



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## Guided Merging of Sequence Diagrams

Why? No more cumbersome manual merging!

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## Formalization

## The tMVML Metamodel



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## Formalization

## State Machine



## Formalization

## State Machine



Given an alphabet $\mathcal{A}_{A}$, a state machine is a quadruple $\left(S, A^{\text {tr }}, A^{\text {eff }}, T\right)$, where

- $S$ is a set of states,
- $A^{t r}, A^{\text {eff }} \subseteq \mathcal{A}_{A}$ are sets of action symbols, and
- $T \subseteq\left(S \times A^{t r} \times \mathcal{P}\left(A^{\text {eff }}\right) \times S\right)$ is a relation representing the transitions between states.


## Formalization

## The tMVML Metamodel



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## Formalization

## Sequence Diagram



## Formalization

## Sequence Diagram



Given the alphabets $\mathcal{A}_{A}$ and $\mathcal{A}_{E}$, and a set $\mathcal{S M}$ of state machines, a sequence diagram is a quadruple ( $L, M$, life, msg), where

- $L$ is a set of lifelines,
- $M$ is a set of messages,
- life : $L \rightarrow\left(\mathcal{S M} \times \mathcal{P}\left(\mathcal{A}_{E}\right) \times \mathcal{P}\left(\mathcal{A}_{E}\right) \times \mathcal{P}\left(\mathcal{A}_{E} \times \mathcal{A}_{E}\right)\right)$
- $\mathrm{msg}: M \rightarrow\left(\mathcal{A}_{A} \times \bigcup_{l \in L} \pi_{2}(\right.$ life $\left.(l)) \times \bigcup_{l \in L} \pi_{3}(\operatorname{life}(l))\right)$.


## Formalization

## Sequence Diagram

Lifeline conformance: For each lifeline, the sequence of received message symbols is a path of triggers in the attached state machine.

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## Sequence Diagram

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## Problem statement

Instance: A sequence diagram and two of its revisions.


## Problem statement

## Objective: Find a consolidated version

- Contains all original and added messages and lifelines
- Lifelines conform to state machines



## Guided Merging of Sequence Diagrams

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## Translation to propositional SAT

Three types of variables

- $m_{i}$, $m$ : message, $i$ : position
- $c_{i}^{s}$, $s$ : state in state machine, $i$ : position, $c$ : "source"
- $t_{i}^{s}, s$ : state in state machine, $i$ : position, $t$ : "target"


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Propositional formula consists of constraints that describe legal solutions, e.g.

$$
\left(m_{1} \vee m_{2} \vee m_{3}\right) \wedge\left(\neg m_{1} \vee \neg m_{2}\right) \wedge\left(\neg m_{2} \vee \neg m_{3}\right) \wedge\left(\neg m_{1} \vee \neg m_{3}\right)
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m_{1} \rightarrow\left(\left(c_{1}^{s 1} \wedge t_{1}^{s 2}\right) \vee\left(c_{1}^{s 3} \wedge t_{1}^{s 4}\right)\right)
\end{gathered}
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\end{gathered}
$$

The encoding is polynomial in the input size. Satisfying assignments of the formula can be directly translated back into a solution of our problem.

## Translation to propositional SAT

$$
\begin{aligned}
& \bigwedge_{m \in M}\left(\bigvee_{i \in \operatorname{allow}(m)} m_{i}\right) \wedge \bigwedge_{\substack{m \in M}}\left(\neg m_{i} \vee \neg m_{j}\right) \\
& \bigwedge_{x \in\{o, \alpha, \beta\}} \bigwedge_{m \in M^{x}}\left(\neg m_{i} \vee \varliminf_{i \in \operatorname{allow}(m)} \varliminf_{\substack { n \in M^{x},{c}{j>i, n \succ m{ n \in M ^ { x } , \begin{subarray} { c } { j > i , \\
n \succ m } }\end{subarray}} n_{j}\right) \\
& \bigwedge_{i \in \operatorname{allow}(m)}\left(\neg m_{i} \vee \bigvee_{t \in \operatorname{trans}(m)}\left(c_{i}^{\pi_{1}(t)} \wedge t_{i}^{\pi_{4}(t)}\right)\right) \\
& \bigwedge_{i=1}^{k}\left(\left(\bigvee_{c_{i}^{s} \in \mathrm{vc}} c_{i}^{s}\right) \wedge\left(\bigvee_{t_{i}^{s} \in \mathrm{vt}}^{s}\right) \wedge \bigwedge_{s \in S_{\text {all }}} \bigwedge_{r \in S_{a l l} \backslash s}\left(\left(\neg c_{i}^{s} \vee \neg c_{i}^{r}\right) \wedge\left(\neg t_{i}^{s} \vee \neg t_{i}^{r}\right)\right)\right) \\
& k-1 \\
& \Lambda \bigwedge \\
& i=1 \quad M \in \mathcal{S} \mathcal{M} s \in \pi_{1}(S M) \\
& \left(\left(t_{i}^{s} \rightarrow \bigwedge_{r \in \pi_{1}(S M) \backslash s} \neg c_{i+1}^{r}\right) \wedge\left(\bigwedge_{j=1}^{i}\left(t_{i}^{s} \wedge \bigwedge_{l=1}^{j} \neg c_{l}^{s} \rightarrow \bigwedge_{r \in \pi_{1}(S M) \backslash s} \neg c_{j+1}^{r}\right)\right)\right)
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## Implementation - Workflow



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## Evaluation

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| Set | \# SM | \# action symbols | \# states | \# transitions |
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| email | 3 | 15 | 16 | 19 |
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## Results

- Between 0.06 s and 0.2 s per solution depending on instance
- Some instances have many solutions $(>1,000)$


## Summary and Future Work

We

- Formalized a subset of the UML in our language $t M V M L$


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- Handling high numbers of solutions


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What we consider next

- Handling high numbers of solutions
- Including deletions and updates
- Integration of other UML concepts
- Visualization


## Related Work

Model merging:

- Gerth et al., Merge support for business process models using term rewriting systems
- Cicchetti et al., Definition of conflict patterns
- Nejati et al., Merging of state machines

Consistency checking:

- Diskin et al., Category theory based framework
- Van der Straeten et al., Inconsistency detection between class and sequence diagrams using Kodkod
- Sabatzadeh et al., Consistency checks between overlapping models
- Tsoliakis, Integration of constraints of other views into sequence diagrams
- Brosch et al., Model checking on state machines and sequence diagrams

