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Model-Driven Engineering



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Model Evolution



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Model Versioning



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Model Versioning

Model Merging



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Model Merging

Multi-View Modelling



























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State Machine



State Machine



Given an alphabet \mathcal{A}_A , a state machine is a quadruple $(S, A^{tr}, A^{e\!f\!f}, T)$, where

- S is a set of states,
- $A^{tr}, A^{e\!f\!f} \subseteq \mathcal{A}_A$ are sets of action symbols, and
- $T \subseteq (S \times A^{tr} \times \mathcal{P}(A^{e\!f\!f}) \times S)$ is a relation representing the transitions between states.


The *tMVML* Metamodel



Sequence Diagram



Sequence Diagram



Given the alphabets \mathcal{A}_A and \mathcal{A}_E , and a set \mathcal{SM} of state machines, a sequence diagram is a quadruple (L, M, life, msg), where

- L is a set of lifelines,
- M is a set of messages,
- $\blacksquare \text{ life}: L \to (\mathcal{SM} \times \mathcal{P}(\mathcal{A}_E) \times \mathcal{P}(\mathcal{A}_E) \times \mathcal{P}(\mathcal{A}_E \times \mathcal{A}_E))$
- msg : $M \to (\mathcal{A}_A \times \bigcup_{l \in L} \pi_2(\mathsf{life}(l)) \times \bigcup_{l \in L} \pi_3(\mathsf{life}(l))).$

Sequence Diagram

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Sequence Diagram



Sequence Diagram



Sequence Diagram



Instance: A sequence diagram and two of its revisions.



Problem statement

Objective: Find a consolidated version

- Contains all original and added messages and lifelines
- Lifelines conform to state machines



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Three types of variables

- m_i , m: message, i: position
- c_i^s , s: state in state machine, i: position, c: "source"
- t_i^s , s: state in state machine, i: position, t: "target"

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Propositional formula consists of constraints that describe legal solutions, e.g.

$$(m_1 \lor m_2 \lor m_3) \land (\neg m_1 \lor \neg m_2) \land (\neg m_2 \lor \neg m_3) \land (\neg m_1 \lor \neg m_3)$$

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The encoding is polynomial in the input size.

Satisfying assignments of the formula can be directly translated back into a solution of our problem.

$$\begin{split} & \bigwedge_{m \in M} \left(\bigvee_{i \in \mathsf{allow}(m)} m_i \right) \wedge \bigwedge_{m \in M} \bigwedge_{i, j \in \mathsf{allow}(m)} \left(\neg m_i \vee \neg m_j \right) \\ & \bigwedge_{x \in \{o, \alpha, \beta\}} \bigwedge_{m \in M^T} \bigwedge_{i \in \mathsf{allow}(m)} \left(\neg m_i \vee \bigvee_{\substack{n \in M^T, \\ n \succ m}} \bigvee_{j \geq i, \\ j \in \mathsf{allow}(n)} n_j \right) \\ & \bigwedge_{i \in \mathsf{allow}(m)} \left(\neg m_i \vee \bigvee_{t \in \mathsf{trans}(m)} \left(c_i^{\pi_1(t)} \wedge t_i^{\pi_4(t)} \right) \right) \\ & \bigwedge_{i = 1}^k \left(\left(\left(\bigvee_{c_i^s \in \mathsf{vc}} c_i^s \right) \wedge \left(\bigvee_{t_i^s \in \mathsf{vt}} t_i^s \right) \wedge \bigwedge_{s \in S_{all}} \bigwedge_{r \in S_{all} \setminus s} \left(\left(\neg c_i^s \vee \neg c_i^r \right) \wedge \left(\neg t_i^s \vee \neg t_i^r \right) \right) \right) \\ & \bigwedge_{i = 1}^{k-1} \bigwedge_{M \in S\mathcal{M}} \bigwedge_{s \in \pi_1(SM)} \left(\left(\bigvee_{j = 1}^i t_i^s \wedge \bigwedge_{j = 1}^j (t_i^s \wedge \bigwedge_{r \in \pi_1(SM) \setminus s} \neg c_{j+1}^r \right) \right) \right) \\ & \left(\left(t_i^s \to \bigwedge_{r \in \pi_1(SM) \setminus s} \neg c_{i+1}^r \right) \wedge \left(\bigwedge_{j = 1}^i \left(t_i^s \wedge \bigwedge_{l = 1}^j \neg c_l^s \to \bigwedge_{r \in \pi_1(SM) \setminus s} \neg c_{j+1}^r \right) \right) \right) \end{split} \right)$$

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Benchmark set with 45 instances

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email	3	15	16	19
coffee	2	9	7	8
philosopher	2	8	7	8

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Results

- Between 0.06s and 0.2s per solution depending on instance
- Some instances have many solutions (>1,000)

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- Including deletions and updates
- Integration of other UML concepts

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Visualization

Related Work

Model merging:

- Gerth et al., Merge support for business process models using term rewriting systems
- *Cicchetti et al.*, Definition of conflict patterns
- Nejati et al., Merging of state machines

Consistency checking:

- Diskin et al., Category theory based framework
- Van der Straeten et al., Inconsistency detection between class and sequence diagrams using Kodkod
- Sabatzadeh et al., Consistency checks between overlapping models
- Tsoliakis, Integration of constraints of other views into sequence diagrams
- Brosch et al., Model checking on state machines and sequence diagrams